

# Single-Deviation Principle and Bargaining

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# Multi-stage games with observable actions

- ▶ finite set of players  $N$
- ▶ stages  $t = 0, 1, 2, \dots$
- ▶  $H$ : set of terminal histories (sequences of action profiles of possibly different lengths)
- ▶ at stage  $t$ , after having observed a non-terminal history of play  $h_t = (a^0, \dots, a^{t-1}) \notin H$ , each player  $i$  simultaneously chooses an action  $a_i^t \in A_i(h_t)$
- ▶  $u_i(h)$ : payoff of  $i \in N$  for terminal history  $h \in H$
- ▶  $\sigma_i$ : behavior strategy for  $i \in N$  specifies  $\sigma_i(h) \in \Delta(A_i(h))$  for  $h \notin H$

Often natural to identify “stages” with time periods.

## Examples

- ▶ repeated games
- ▶ alternating bargaining game

# Unimprovable Strategies

To verify that a strategy profile  $\sigma$  constitutes a subgame perfect equilibrium (SPE) in a multi-stage game with observed actions, it suffices to check whether there are any histories  $h_t$  where some player  $i$  can gain by deviating from playing  $\sigma_i(h_t)$  at  $t$  and conforming to  $\sigma_i$  elsewhere.

$u_i(\sigma|h_t)$ : expected payoff of player  $i$  in the subgame starting at  $h_t$  and played according to  $\sigma$  thereafter

## Definition 1

A strategy  $\sigma_i$  is *unimprovable* given  $\sigma_{-i}$  if  $u_i(\sigma_i, \sigma_{-i} | h_t) \geq u_i(\sigma'_i, \sigma_{-i} | h_t)$  for every  $h_t$  and  $\sigma'_i$  with  $\sigma'_i(h) = \sigma_i(h)$  for all  $h \neq h_t$ .

# Continuity at Infinity

If  $\sigma$  is an SPE then  $\sigma_i$  is unimprovable given  $\sigma_{-i}$ . For the converse...

## Definition 2

A game is *continuous at infinity* if

$$\lim_{t \rightarrow \infty} \sup_{\{(h, \tilde{h}) | h_t = \tilde{h}_t\}} |u_i(h) - u_i(\tilde{h})| = 0, \forall i \in N.$$

Events in the distant future are relatively unimportant.

# Single (or One-Shot) Deviation Principle

## Theorem 1

Consider a multi-stage game with observed actions that is *continuous at infinity*. If  $\sigma_i$  is *unimprovable* given  $\sigma_{-i}$  for all  $i \in N$ , then  $\sigma$  constitutes an SPE.

Proof allows for infinite action spaces at some stages. There exist versions for games with unobserved actions.

## Proof

Suppose that  $\sigma_i$  is unimprovable given  $\sigma_{-i}$ , but  $\sigma_i$  is not a best response to  $\sigma_{-i}$  following some history  $h_t$ . Let  $\sigma_i^1$  be a strictly better response and define

$$\varepsilon = u_i(\sigma_i^1, \sigma_{-i}|h_t) - u_i(\sigma_i, \sigma_{-i}|h_t) > 0.$$

Since the game is *continuous at infinity*, there exists  $t' > t$  and  $\sigma_i^2$  s.t.  $\sigma_i^2$  is identical to  $\sigma_i^1$  at all information sets up to (and including) stage  $t'$ ,  $\sigma_i^2$  coincides with  $\sigma_i$  across all longer histories and

$$|u_i(\sigma_i^2, \sigma_{-i}|h_t) - u_i(\sigma_i^1, \sigma_{-i}|h_t)| < \varepsilon/2.$$

Then

$$u_i(\sigma_i^2, \sigma_{-i}|h_t) > u_i(\sigma_i, \sigma_{-i}|h_t).$$

## Proof

$\sigma_i^3$ : strategy obtained from  $\sigma_i^2$  by replacing the stage  $t'$  actions following any history  $h_{t'}$  with the corresponding actions under  $\sigma_i$

Conditional on any  $h_{t'}$ ,  $\sigma_i$  and  $\sigma_i^3$  coincide, hence

$$u_i(\sigma_i^3, \sigma_{-i} | h_{t'}) = u_i(\sigma_i, \sigma_{-i} | h_{t'}).$$

As  $\sigma_i$  is *unimprovable* given  $\sigma_{-i}$ , and conditional on  $h_{t'}$  the subsequent play in strategies  $\sigma_i$  and  $\sigma_i^2$  differs only at stage  $t'$ ,

$$u_i(\sigma_i, \sigma_{-i} | h_{t'}) \geq u_i(\sigma_i^2, \sigma_{-i} | h_{t'}).$$

Then

$$u_i(\sigma_i^3, \sigma_{-i} | h_{t'}) \geq u_i(\sigma_i^2, \sigma_{-i} | h_{t'})$$

for all histories  $h_{t'}$ . Since  $\sigma_i^2$  and  $\sigma_i^3$  coincide before reaching stage  $t'$ ,

$$u_i(\sigma_i^3, \sigma_{-i} | h_t) \geq u_i(\sigma_i^2, \sigma_{-i} | h_t).$$

# Proof

$\sigma_i^4$ : strategy obtained from  $\sigma_i^3$  by replacing the stage  $t' - 1$  actions following any history  $h_{t'-1}$  with the corresponding actions under  $\sigma_i$

Similarly,

$$u_i(\sigma_i^4, \sigma_{-i}|h_t) \geq u_i(\sigma_i^3, \sigma_{-i}|h_t) \dots$$

The final strategy  $\sigma_i^{t'-t+3}$  is identical to  $\sigma_i$  conditional on  $h_t$  and

$$\begin{aligned} u_i(\sigma_i, \sigma_{-i}|h_t) &= u_i(\sigma_i^{t'-t+3}, \sigma_{-i}|h_t) \geq \dots \\ &\geq u_i(\sigma_i^3, \sigma_{-i}|h_t) \geq u_i(\sigma_i^2, \sigma_{-i}|h_t) > u_i(\sigma_i, \sigma_{-i}|h_t), \end{aligned}$$

a contradiction.

# Applications

Apply the single deviation principle to repeated prisoners' dilemma to implement the following equilibrium paths for high discount factors:

- ▶  $(C, C), (C, C), \dots$
- ▶  $(C, C), (C, C), (D, D), (C, C), (C, C), (D, D), \dots$
- ▶  $(C, D), (D, C), (C, D), (D, C) \dots$

	$C$	$D$
$C$	1, 1	-1, 2
$D$	2, -1	0, 0

Cooperation is possible in repeated play.

# Bargaining with Alternating Offers

Rubinstein (1982)

- ▶ players  $i = 1, 2; j = 3 - i$
- ▶ set of feasible utility pairs

$$U = \{(u_1, u_2) \in [0, \infty)^2 \mid u_2 \leq g_2(u_1)\}$$

- ▶  $g_2$  s. decreasing, concave (and hence continuous),  $g_2(0) > 0$
- ▶  $\delta_i$ : discount factor of player  $i$
- ▶ at every time  $t = 0, 1, \dots$ , player  $i(t)$  proposes an alternative  $u = (u_1, u_2) \in U$  to player  $j(t) = 3 - i(t)$

$$i(t) = \begin{cases} 1 & \text{for } t \text{ even} \\ 2 & \text{for } t \text{ odd} \end{cases}$$

- ▶ if  $j(t)$  accepts the offer, game ends yielding payoffs  $(\delta_1^t u_1, \delta_2^t u_2)$
- ▶ otherwise, game proceeds to period  $t + 1$

# Stationary SPE

Define  $g_1 = g_2^{-1}$ . Graphs of  $g_2$  and  $g_1^{-1}$ : Pareto-frontier of  $U$

Let  $(m_1, m_2)$  be the unique solution to the following system of equations

$$\begin{aligned}m_1 &= \delta_1 g_1(m_2) \\m_2 &= \delta_2 g_2(m_1).\end{aligned}$$

$(m_1, m_2)$  is the intersection of the graphs of  $\delta_2 g_2$  and  $(\delta_1 g_1)^{-1}$ .

SPE in “stationary” strategies: in any period where player  $i$  has to make an offer to  $j$ , he offers  $u$  with  $u_j = m_j$  and  $u_i = g_i(m_j)$ , and  $j$  accepts only offers  $u$  with  $u_j \geq m_j$ .

*Single-deviation principle*: constructed strategies form an SPE.

Is the SPE unique?

# Iterated Conditional Dominance

## Definition 3

In a multi-stage game with observable actions, an action  $a_i$  is *conditionally dominated* at stage  $t$  given history  $h_t$  if, in the subgame starting at  $h_t$ , every strategy for player  $i$  that assigns positive probability to  $a_i$  is strictly dominated.

## Proposition 1

*In any multi-stage game with observable actions, every SPE survives the iterated elimination of conditionally dominated strategies.*

# Equilibrium uniqueness

*Iterated conditional dominance*: stationary equilibrium is essentially the unique SPE.

## Theorem 2

*The SPE of the alternating-offer bargaining game is unique, except for the decision to accept or reject Pareto-inefficient offers.*

# Proof

- ▶ Following a disagreement at date  $t$ , player  $i$  cannot obtain a period  $t$  expected payoff greater than

$$M_i^0 = \delta_i \max_{u \in U} u_i = \delta_i g_i(0)$$

- ▶ Rejecting an offer  $u$  with  $u_i > M_i^0$  is conditionally dominated by accepting such an offer for  $i$ .
- ▶ Once we eliminate dominated actions,  $i$  accepts all offers  $u$  with  $u_i > M_i^0$  from  $j$ .
- ▶ Making any offer  $u$  with  $u_i > M_i^0$  is dominated for  $j$  by an offer  $\bar{u} = \lambda u + (1 - \lambda) (M_i^0, g_j(M_i^0))$  for  $\lambda \in (0, 1)$  (both offers are accepted immediately).

# Proof

Under the surviving strategies

- ▶  $j$  can reject an offer from  $i$  and make a counteroffer next period that leaves him with slightly less than  $g_j(M_i^0)$ , which  $i$  accepts; it is conditionally dominated for  $j$  to accept any offer smaller than

$$m_j^1 = \delta_j g_j(M_i^0)$$

- ▶  $i$  cannot expect to receive a continuation payoff greater than

$$M_i^1 = \max(\delta_i g_i(m_j^1), \delta_i^2 M_i^0) = \delta_i g_i(m_j^1)$$

after rejecting an offer from  $j$

$$\delta_i g_i(m_j^1) = \delta_i g_i(\delta_j g_j(M_i^0)) \geq \delta_i g_i(g_j(M_i^0)) = \delta_i M_i^0 \geq \delta_i^2 M_i^0$$

# Proof

Recursively define

$$\begin{aligned}m_j^{k+1} &= \delta_j g_j(M_i^k) \\ M_i^{k+1} &= \delta_i g_i(m_j^{k+1})\end{aligned}$$

for  $i = 1, 2$  and  $k \geq 1$ .  $(m_i^k)_{k \geq 0}$  is increasing and  $(M_i^k)_{k \geq 0}$  is decreasing.

Prove by induction on  $k$  that, under any strategy that survives iterated conditional dominance, player  $i = 1, 2$

- ▶ never accepts offers with  $u_i < m_i^k$
- ▶ always accepts offers with  $u_i > M_i^k$ , but making such offers is dominated for  $j$ .

# Proof

- ▶ The sequences  $(m_i^k)$  and  $(M_i^k)$  are monotonic and bounded, so they need to converge. The limits satisfy

$$\begin{aligned}m_j^\infty &= \delta_j g_j(\delta_i g_i(m_j^\infty)) \\M_i^\infty &= \delta_i g_i(m_j^\infty).\end{aligned}$$

- ▶  $(m_1^\infty, m_2^\infty)$  is the (unique) intersection point of the graphs of the functions  $\delta_2 g_2$  and  $(\delta_1 g_1)^{-1}$
- ▶  $M_i^\infty = \delta_i g_i(m_j^\infty) = m_i^\infty$
- ▶ All strategies of  $i$  that survive iterated conditional dominance accept  $u$  with  $u_i > M_i^\infty = m_i^\infty$  and reject  $u$  with  $u_i < m_i^\infty = M_i^\infty$ .

# Proof

In an SPE

- ▶ at any history where  $i$  is the proposer,  $i$ 's payoff is at least  $g_i(m_j^\infty)$ : offer  $u$  arbitrarily close to  $(g_i(m_j^\infty), m_j^\infty)$ , which  $j$  accepts under the strategies surviving the elimination process
- ▶  $i$  cannot get more than  $g_i(m_j^\infty)$ 
  - ▶ any offer made by  $i$  specifying a payoff greater than  $g_i(m_j^\infty)$  for himself would leave  $j$  with less than  $m_j^\infty$ ; such offers are rejected by  $j$  under the surviving strategies
  - ▶ under the surviving strategies,  $j$  never offers  $i$  more than  $M_i^\infty = \delta_i g_i(m_j^\infty) \leq g_i(m_j^\infty)$
- ▶ hence  $i$ 's payoff at any history where  $i$  is the proposer is exactly  $g_i(m_j^\infty)$ ; possible only if  $i$  offers  $(g_i(m_j^\infty), m_j^\infty)$  and  $j$  accepts with **probability 1**

Uniquely pinned down actions at every history, except those where  $j$  has just received an offer  $(u_i, m_j^\infty)$  for some  $u_i < g_i(m_j^\infty)$ ...

# Properties of the equilibrium

- ▶ The SPE is **efficient**—agreement is obtained in the first period, without delay.
- ▶ SPE payoffs:  $(g_1(m_2), m_2)$ , where  $(m_1, m_2)$  solve

$$m_1 = \delta_1 g_1(m_2)$$

$$m_2 = \delta_2 g_2(m_1).$$

- ▶ **Patient** players get higher payoffs: the payoff of player  $i$  is increasing in  $\delta_i$  and decreasing in  $\delta_j$ .
- ▶ For a fixed  $\delta_1 \in (0, 1)$ , the payoff of player 2 converges to 0 as  $\delta_2 \rightarrow 0$  and to  $\max_{u \in U} u_2$  as  $\delta_2 \rightarrow 1$ .
- ▶ If  $U$  is symmetric and  $\delta_1 = \delta_2$ , player 1 enjoys a **first mover advantage**:  $m_1 = m_2$  and  $g_1(m_2) = m_2/\delta > m_2$ .

# Nash Bargaining

Assume  $g_2$  is decreasing, s. concave and continuously differentiable.

*Nash (1950) bargaining solution:*

$$\{u^*\} = \arg \max_{u \in U} u_1 u_2 = \arg \max_{u \in U} u_1 g_2(u_1).$$

## Theorem 3 (Binmore, Rubinstein and Wolinsky 1985)

*Suppose that  $\delta_1 = \delta_2 =: \delta$  in the alternating bargaining model. Then the unique SPE payoffs converge to the Nash bargaining solution as  $\delta \rightarrow 1$ .*

$$m_1 g_2(m_1) = m_2 g_1(m_2)$$

$(m_1, g_2(m_1))$  and  $(g_1(m_2), m_2)$  belong to the intersection of  $g_2$ 's graph with the same hyperbola, which approaches the hyperbola tangent to the boundary of  $U$  (at  $u^*$ ) as  $\delta \rightarrow 1$ .

# Bargaining with random selection of proposer

- ▶ Two players need to divide \$1.
- ▶ Every period  $t = 0, 1, \dots$  player 1 is chosen with probability  $p$  to make an offer to player 2.
- ▶ Player 2 accepts or rejects 1's proposal.
- ▶ Roles are interchanged with probability  $1 - p$ .
- ▶ In case of disagreement the game proceeds to the next period.
- ▶ The game ends as soon as an offer is accepted.
- ▶ Player  $i = 1, 2$  has discount factor  $\delta_i$ .

# Equilibrium

- ▶ The unique equilibrium is **stationary**, i.e., each player  $i$  has the same expected payoff  $v_i$  in every subgame.
- ▶ Payoffs solve

$$\begin{aligned}v_1 &= p(1 - \delta_2 v_2) + (1 - p)\delta_1 v_1 \\v_2 &= p\delta_2 v_2 + (1 - p)(1 - \delta_1 v_1).\end{aligned}$$

- ▶ The solution is

$$\begin{aligned}v_1 &= \frac{p/(1 - \delta_1)}{p/(1 - \delta_1) + (1 - p)/(1 - \delta_2)} \\v_2 &= \frac{(1 - p)/(1 - \delta_2)}{p/(1 - \delta_1) + (1 - p)/(1 - \delta_2)}.\end{aligned}$$

# Comparative Statics

$$v_1 = \frac{1}{1 + \frac{(1-p)(1-\delta_1)}{p(1-\delta_2)}}$$
$$v_2 = \frac{1}{1 + \frac{p(1-\delta_2)}{(1-p)(1-\delta_1)}}.$$

- ▶ **Immediate agreement**
- ▶ **First mover advantage**
  - ▶  $v_1$  increases with  $p$ ,  $v_2$  decreases with  $p$ .
  - ▶ For  $\delta_1 = \delta_2$ , we obtain  $v_1 = p$ ,  $v_2 = 1 - p$ .
- ▶ **Patience pays off**
  - ▶  $v_i$  increases with  $\delta_i$  and decreases with  $\delta_j$  ( $j = 3 - i$ ).
  - ▶ Fix  $\delta_j$  and take  $\delta_i \rightarrow 1$ , we get  $v_i \rightarrow 1$  and  $v_j \rightarrow 0$ .

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