

Second proof of remark given last time:

Assume $x_0 \subseteq x_1 \subseteq x_2 \subseteq \dots$ are increasing tuples of variables.
 ~~$p_0(x_0) \subseteq p_1(x_1) \subseteq \dots$~~
 $p_0(x_0) \subseteq p_1(x_1) \subseteq \dots$ are increasing complete types

Let $a_i \models p_i \quad \forall i$

~~For each $i \leq j$, let $a_{j,i}$ be the subtuple of a_j corresponding to x_i .~~

For each $i \leq j$, let $a_{j,i}$ be the subtuple of a_j corresponding to x_i .

Define $b_0 \subseteq b_1 \subseteq b_2 \subseteq \dots$ st $b_i \models p_i$

Let $b_0 = a_0$. Assuming we have b_i , the $\text{tp}(b_i) = \text{tp}(a_{i+1,i}) = p_i$

so there is an automorphism f sending $a_{i+1,i}$ to b_i .

Let $b_{i+1} = f(a_{i+1})$. Let $b = \bigcup b_i$. Then $b \models \bigcup p_i$.

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Definition A partial type $p(x, b)$ divides over c

(x, b, c) are possibly infinite tuples of variables/elements of \mathcal{U}

if there is an indiscernible sequence $(b_i : i < \omega)$ in $tp(b/c)$

st. $\bigwedge_{i < \omega} p(x, b_i)$ [= $U p(x, b_i)$] is inconsistent.

Remark: If $(b_i : i < \omega)$ is an indiscernible sequence in $tp(b/c)$ then it has an automorphic image which is also c -indiscernible.

Proof: ~~There is~~

By compactness, for every λ , there is a similar sequence $\neq \emptyset$ in $tp(b/c)$. [Let $p_n(x_0, \dots, x_{n-1}) = tp(b_0, \dots, b_{n-1}) = \bigwedge_{i < \lambda} q_i(x_i) \wedge$

$\bigwedge_{i_0 < \dots < i_{n-1}} p_n(x_{i_0}, \dots, x_{i_{n-1}})$
consistent.]

Now extract a c -indiscernible sequence: $(b_i'' : i < \omega)$
 \leftarrow Is b_i' the similar sequence in $tp(b/c)$ of length λ ? Yes

① $b_i' \models q_i$ for all $i \Rightarrow b_i'' \models q_i$

② $\forall i_0 < \dots < i_{n-1} \models p_n(b_{i_0}', \dots, b_{i_{n-1}}') \Rightarrow \models p_n(b_{i_0}'', \dots, b_{i_{n-1}}'')$

$\Rightarrow tp(b_i : i < \omega) = tp(b_i'' : i < \omega)$.

$\Rightarrow b_{<\omega}''$ is an automorphic image of $b_{<\omega}$. \square

Definition $a \not\downarrow_c b$ (read: a independent of b over c)

if $tp(a/bc)$ does not divide/c.
and

Proposition. $a \downarrow_c b$ if and only if every indiscernible sequence in ${}^c_t p(b/c)$ has an c -automorphic image in $tp(b/ac)$.

Proof: \Rightarrow : Assume $a \not\downarrow_c b$ i.e. $tp(a/bc)$ does not divide (dnd) over c .

write $p(x, bc) = tp(a/bc)$.

Let $(b_i : i < \omega)$ be an c -indiscernible sequence in $tp(b/c)$.

~~Then $(b_i c : i < \omega)$ is indiscernible in $tp(b/c/c)$.~~

[By the remark, there is an automorphic image $(b_i' : i < \omega)$ which is c -indiscernible and in $tp(b/c)$.]

$\Rightarrow (b_i' c : i < \omega)$ is ~~an~~ indiscernible in $tp(b/c/c)$.

Since $tp(a/bc)$ dnd over c , there is $a' \neq \bigwedge p(x, b_i' c)$.

In particular, $a' \neq tp(a/c)$.

~~Mapping~~ Let f be an c -automorphism st. $f(a/c) = ac$.

Therefore $(b_i'') := f(b_i')$ is an c -automorphic image of (b_i)

and $\bigwedge p(a, b_i'' c) \Rightarrow b_i'' \neq tp(b/ac)$.

\Leftarrow : Let $(b_i c)$ be any ~~an~~ c -indiscernible sequence in $tp(b/c/c) \Rightarrow \frac{a=c}{\forall i}$

~~is indiscernible~~

We need to find $a' \neq \bigwedge p(x, b_i c)$.

By assumption (b_i) has a c -automorphic image (b_i') in $tp(b/ac)$.

Let f be the c -automorphism & let $a' = f^{-1}(a)$.

Then $\bigwedge p(a, b/c) \Rightarrow \bigwedge p(a', b/c)$. □

Corollary: (1) Downward right-hand transitivity:

$$\forall a, b, c, d: \quad a \downarrow_c b, d \Rightarrow a \downarrow_c b \wedge a \downarrow_{bc} d.$$

(2) Upward left-hand transitivity:

$$a \downarrow_c b \text{ and } d \downarrow_{ac} b \Rightarrow ad \downarrow_c b.$$

Proof (1) Assume $a \downarrow_c b, d$ then ~~if~~ if (a) is c -indiscernible in $tp(b/c)$ then by extension/extraction, we can find (d_i) st.

$(b_i d_i)$ is c -indiscernible in $tp(b, d/c)$.

(Extend to $(b_i : i < \lambda)$, for each i find d_i st. $b_i d_i \equiv_c b, d$
(by not sending b to b_i)

and extract a c -indiscernible sequence $(b_i' d_i')$.

but $b \prec_w \equiv_c b' \prec_w$ (both have similar ~~sequences~~ ^{c -indiscernible} seq of same length.

so we may _(w.m.a) assume $b \prec_w = b' \prec_w$.

Since $a \downarrow_c b, d$, there is a c -automorphic image $(b_i' d_i')$ in $tp(b, d/c)$

in particular (b_i') is a c -automorphic image of (b_i) in

$$tp(b/c) \Rightarrow a \downarrow_c b.$$

Now: let $(b_i d_i)$ be b_i -indiscernible in $tp(b, d/c)$.

Then ~~(b_i d_i)~~ ^{it} is c -indiscernible in $tp(b, d/c)$.

$$\text{Let } tp(a/bcd) = q(x/bcd).$$

So $\bigwedge q(x, bc, d_i)$ is consistent $\Rightarrow a \downarrow_{bc} d$.

(2) Assume $a \downarrow_c b, d \downarrow_{ac} b$.

Let (b_i) be a c -indiscernible sequence in $tp(b/c)$.

Since $a \downarrow_c b$ there is a c -automorphic image (b_i') in $tp(b/ac)$.

~~By an earlier remark~~

b_i' is c -indiscernible and by a previous remark has c -automorphic image which is still in $tp(b/ac)$ and in addition is ac -indiscernible.

So we may assume (b_i') is ac -indiscernible.

Since $d \downarrow_{ac} b$, then (b_i') has an ac -automorphic image in $tp(b/acd)$.

Conclusion (b_i) has a c -automorphic image in $tp(b/acd)$.

$\Rightarrow a d \downarrow_c b$

□

Lemma A partial type $p(x, b)$ divides $/c$ iff there is a formula $\varphi(x, b) \in p(x, b)$ which does.

(convention: all partial types are closed under conjunction)

Proof \Leftarrow : clear.

\Rightarrow : Assume (b_i) is c -indiscernible and $\bigcup p(x, b_i)$ is inconsistent. By compactness, only finitely many

formulas are required for inconsistency, say

$$\varphi_0(x, b_{i_0}) \in p(x, b_{i_0}), \dots, \varphi_{k-1}(x, b_{i_{k-1}}) \in p(x, b_{i_{k-1}}).$$

let $\psi = \bigwedge \varphi_i(x, y)$. Then $\psi(x, b) \in p(x, b)$.

and $\bigwedge \psi(x, b_i)$ is inconsistent $\Rightarrow \psi(x, b)$ divides $/c$. \square

Corollary Finite Character $a \not\downarrow_c b$ ~~iff~~ $(a, b, c$ are possibly infinite)
 $\Leftrightarrow \forall a' \subseteq a$ and $\forall b' \subseteq b$ finite, $a' \not\downarrow_c b'$

Proof \Rightarrow : clear.

\Leftarrow : if $a \not\downarrow_c b$ then there is a formula $\varphi(x, bc) \in tp(a/c)$ which divides over c .

now only finite subtuples $a' \subseteq a$ and $b' \subseteq b$ actually appear in φ . let $a' \subseteq a$ correspond to $x' \subseteq x$.

$$\Rightarrow \varphi(x', b'c) \in tp(a'/b'c).$$

$$\Rightarrow a' \not\downarrow_c b'$$

\square .