

18.S096: Homework Problem Set 2

Topics in Mathematics of Data Science (Fall 2015)

Afonso S. Bandeira

Due on October 20, 2015

Problem 2.1 Given a graph $G = (V, E, W)$ consider the random walk on V with transition probabilities

$$M_{ij} = \text{Prob} \{X(t+1) = j | X(t) = i\} = \frac{w_{ij}}{\text{deg}(i)}.$$

Partition the vertex set as $V = V_+ \cup V_- \cup V_*$. Suppose that every node in V_* is connected to at least a node in either V_+ or V_- . Given a node $i \in V$ let $g(i)$ be the probability that a random walker starting at i reaches a node in V_+ before reaching one in V_- , i.e.:

$$g(i) = \text{Prob} \left\{ \inf_{t \geq 0: X(t) \in V_+} t < \inf_{t \geq 0: X(t) \in V_-} t | X(0) = i \right\}.$$

Note that if $i \in V_+$ then $g(i) = 1$ and, if $i \in V_-$, then $g(i) = 0$. What is the value of g in V_* ? How would you compute it?

Problem 2.2 For a graph G let $h(G)$ denote its Cheeger constant and $\lambda_2(\mathcal{L}_G)$ the second smallest eigenvalue of its normalized graph Laplacian. Recall that Cheeger inequality guarantees that

$$\frac{1}{2} \lambda_2(\mathcal{L}_G) \leq h_G \leq \sqrt{2 \lambda_2(\mathcal{L}_G)}.$$

This exercise shows that this inequality is tight (at least up to constants).

1. Construct a family of graphs for which $\lambda_2(\mathcal{L}_G) \rightarrow 0$ and for which there exists a constant $C > 0$ for which, for every G in the family,

$$h_G \leq C \lambda_2(\mathcal{L}_G)$$

2. Construct a family of graphs for which $\lambda_2(\mathcal{L}_G) \rightarrow 0$ and for which there exists a constant $c > 0$ for which, for every G in the family

$$h_G \geq c\sqrt{\lambda_2(\mathcal{L}_G)}$$

Problem 2.3 Given a graph G show that the dimension of the nullspace of L_G corresponds to the number of connected components of G .

Problem 2.4 Given a connected unweighted graph $G = (V, E)$, its diameter is equal to

$$\text{diam}(G) = \max_{u,v \in V} \min_{\text{path } p \text{ from } u \text{ to } v} \text{length of } p.$$

Show that

$$\text{diam}(G) \geq \frac{1}{\text{vol}(G)\lambda_2(\mathcal{L}_G)}.$$

Problem 2.5 1. Prove the Courant Fisher Theorem: Given a symmetric matrix $A \in \mathbb{R}^{n \times n}$ with eigenvalues $\lambda_1 \leq \dots \leq \lambda_n$, for $k \leq n$,

$$\lambda_k(A) = \min_{U: \dim(U)=k} \left[\max_{x \in U} \frac{x^T A x}{x^T x} \right].$$

2. Show also that:

$$\lambda_2(A) = \max_{y \in \mathbb{R}^n} \left[\min_{x \in \mathbb{R}^n: x \perp y} \frac{x^T A x}{x^T x} \right].$$

Problem 2.6 Given a set of points $x_1, \dots, x_n \in \mathbb{R}^p$ and a partition of them in k clusters S_1, \dots, S_k recall the k -means objective

$$\min_{S_1, \dots, S_k} \min_{\mu_1, \dots, \mu_k} \sum_{l=1}^k \sum_{i \in S_l} \|x_i - \mu_l\|^2.$$

Show that this is equivalent to

$$\min_{S_1, \dots, S_k} \sum_{l=1}^k \frac{1}{|S_l|} \sum_{i,j \in S_l} \|x_i - x_j\|^2.$$

MIT OpenCourseWare
<http://ocw.mit.edu>

18.S096 Topics in Mathematics of Data Science
Fall 2015

For information about citing these materials or our Terms of Use, visit: <http://ocw.mit.edu/terms>.