

# **Chapter 19. Meeting 19, Approaches: Grammars and L-Systems**

## **19.1. Announcements**

- Sonic system draft due: 27 April
- No class Tuesday, 20 April
- Be sure to do reading for next class:

Riskin, J. 2003. "The Defecating Duck, or, the Ambiguous Origins of Artificial Life." *Critical Inquiry* 29(4): 599-633.

## **19.2. Quiz**

- 10 Minutes

## **19.3. String Rewriting Systems**

- Given an alphabet and rewrite (production) rules, transform strings
- A wide variety of formalizations and approaches
- Axel Thue: first systematic treatment
- Noam Chomsky: applied concept of re-writing to syntax of natural languages

## **19.4. Formal Grammars**

- A set of rules for a formal language
- Formal grammars can be generative or analytic
- Generative grammars defined by
  - A finite set of nonterminal symbols (variables that can be replaced)
  - A finite set of terminal symbols (constants)
  - An axiom, or initial state
  - A finite set of production rules, replacing variables with variables or constants
- Generative grammars are iterative

## 19.5. Lindenmayer Systems

- Based on 1968 work of Aristid Lindenmayer
- Origins in model of a natural systems: “a theoretical framework for studying the development of simple multicellular organisms”
- 1984: began use of using computer graphics for visualization of plan structures
- L-systems: formal grammars where re-writing is parallel, not sequential: all symbols are simultaneously replaced

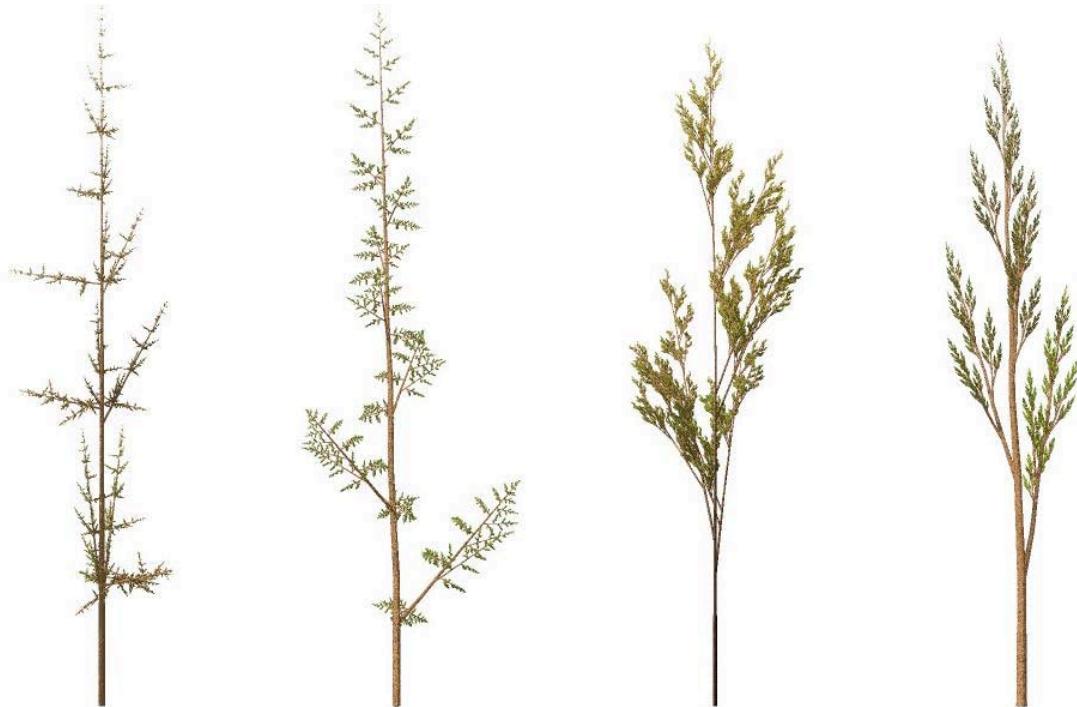


Image: Public domain (Wikipedia)

YouTube (<http://www.youtube.com/watch?v=L54SE9KTMSQ>)

YouTube (<http://www.youtube.com/watch?v=t-FZhw9G-RQ>)

YouTube (<http://www.youtube.com/watch?v=t-FZhw9G-RQ>)

- Motivation from natural systems: idea of cell divisions of occurring at the same time
- L-systems can take many different forms depending on rule systems and alphabet components

## 19.6. Context-Free

- Rules match one source to one or more destination
- Example:

**Alphabet:**

**V: A B**

**Production rules :**

P1:A → AB

P2:B → A

**axiom:**

$\omega : B$

which produces for derivation step  $n$ :

$n=0 : B$

$n=1 : A$

$n=2 : AB$

$n=3 : ABA$

$n=4 : ABAAB$

$n=5 : ABAABABA$

Courtesy of Stelios Manousakis. Used with permission. From "Musical L-Systems." Master's Thesis, Royal Conservatory, The Hague, 2006.

- Originally proposed by Lindenmayer to model growth of algae
- Graphic representation Prusinkiewicz and Lindenmayer (1990)



Figure 1.3: Example of a derivation in a DOL-system.

© P. Prusinkiewicz and A. Lindenmayer (from *The Algorithmic Beauty of Plants*).  
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## 19.7. Context-Sensitive

- Rules match two or more sources to one or more destination
- 1L systems: match left or right of target source
- 2L systems: match left and right of target source
- 1L systems can be considered 2L systems with an empty (open matching) context
- Example:

**Alphabet:**

$V:ab$

**Production rules:**

P1: $b < a > \emptyset \rightarrow b$

P2: $b \rightarrow a$

**axiom:**

$\omega : baaaaaaaaaa$

which produces for derivation step  $n$ :

$n=0 : baaaaaaaaaa$

$n=1 : abaaaaaaaaaa$

$n=2 : aabaaaaaaaaaa$

$n=3 : aaabaaaaaaaaaa$

$n=4 : aaaabaaaaaaaaaa$

Courtesy of Stelios Manousakis. Used with permission. From "Musical L-Systems." Master's Thesis, Royal Conservatory, The Hague, 2006.

## 19.8. Non-Deterministic and Table L-systems

- A context-sensitive or context free grammar can be deterministic
- If the application of rules is probabilistic, non-deterministic grammar is created
- Common approach: map one source to two or more destinations, with weighted probabilities for each destination
- Example:

**Alphabet :**

**V: A B**

**Production rules :**

P1: A  $\xrightarrow{70\%}$  AB

P2: A  $\xrightarrow{30\%}$  BA

P3: B  $\longrightarrow$  A

**axiom :**

$\omega : A$

which can produce for derivation step n:

n=0 : A

n=1 : AB

n=2 : ABA

n=3 : BAAAB

n=4 : ABAABBAA

or:

n=0 : A

n=1 : BA

n=2 : AAB

n=3 : ABABA

n=4 : BAABAAAB

Courtesy of Stelios Manousakis. Used with permission. From "Musical L-Systems." Master's Thesis, Royal Conservatory, The Hague, 2006.

- Alternatively, rules can be changed during production, producing a Table L-system (Manousakis 2006, p. 29)

- Example:

**Alphabet:**

**V: A B**

**axiom:**

**$\omega: B$**

**Table 1:**

**Production rules :**

**$P1: A \rightarrow AB$**

**$P2: B \rightarrow A$**

**Table 2:**

**Production rules :**

**$P1: A \rightarrow B$**

**$P2: B \rightarrow BA$**

If the set changes on derivation step n=3, this would produce:

<b>T1</b>	<b>n=0 : B</b>
	<b>n=1 : A</b>
	<b>n=2 : AB</b>
<b>T2</b>	<b>n=3 : BBA</b>
	<b>n=4 : BABAB</b>
	<b>n=5 : BABBABBA</b>

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## 19.9. Non-Propagative L-systems

- Where rules replace source with more than one successor, the system grows and is propagative
- If rules only encode one-value destinations, the rule system is non-propagative

- Context-sensitive non-propagative L-systems are identical to a standard 1D CA
- Example:

**Alphabet :**

V : A B

**Production rules :**

P1:A < A > A → B

P2:A < A > B → A

P3:A < B > A → B

P4:A < B > B → A

P5:B < A > A → A

P6:B < A > B → B

P7:B < B > A → A

P8:B < B > B → B

**axiom :**

$\omega$ : BBBB BBBB BBBB BBBB BBBB

Using the classic CA visualization for this grammar, and interpreting A = 1 (black pixel) and B = 0 (grey pixel), the first 30 generations look like this:



Figure 2.3. Cellular automata with L-systems.

Courtesy of Stelios Manousakis. Used with permission. From "Musical L-Systems." Master's Thesis, Royal Conservatory, The Hague, 2006.

## 19.10. Musical and Artistic Application of L-systems

- First published implementation: Prusinkiewicz

Prusinkiewicz, P. 1986. "Score Generation with L-Systems." In *Proceedings of the International Computer Music Conference*. San Francisco: International Computer Music Association. 455-457.

- A spatial mapping of 2D graphical output of L-system curves to pitch (vertical) and duration (horizontal)

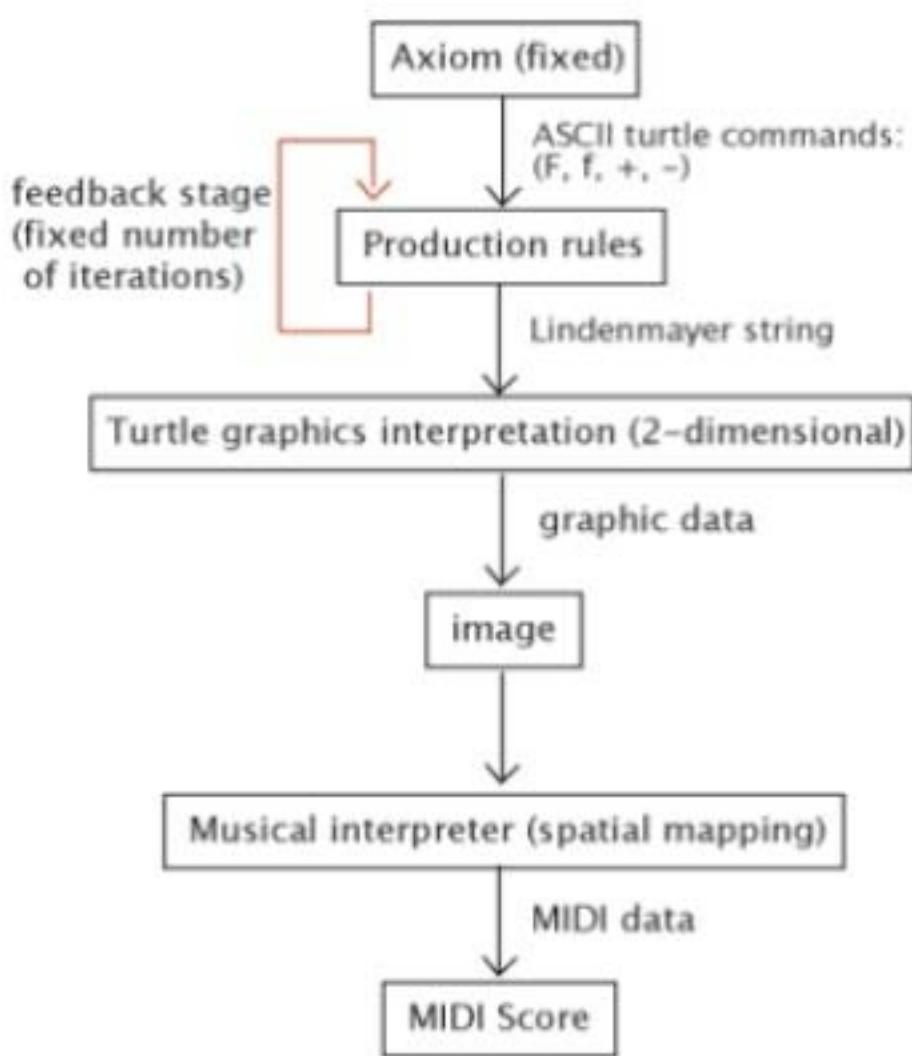
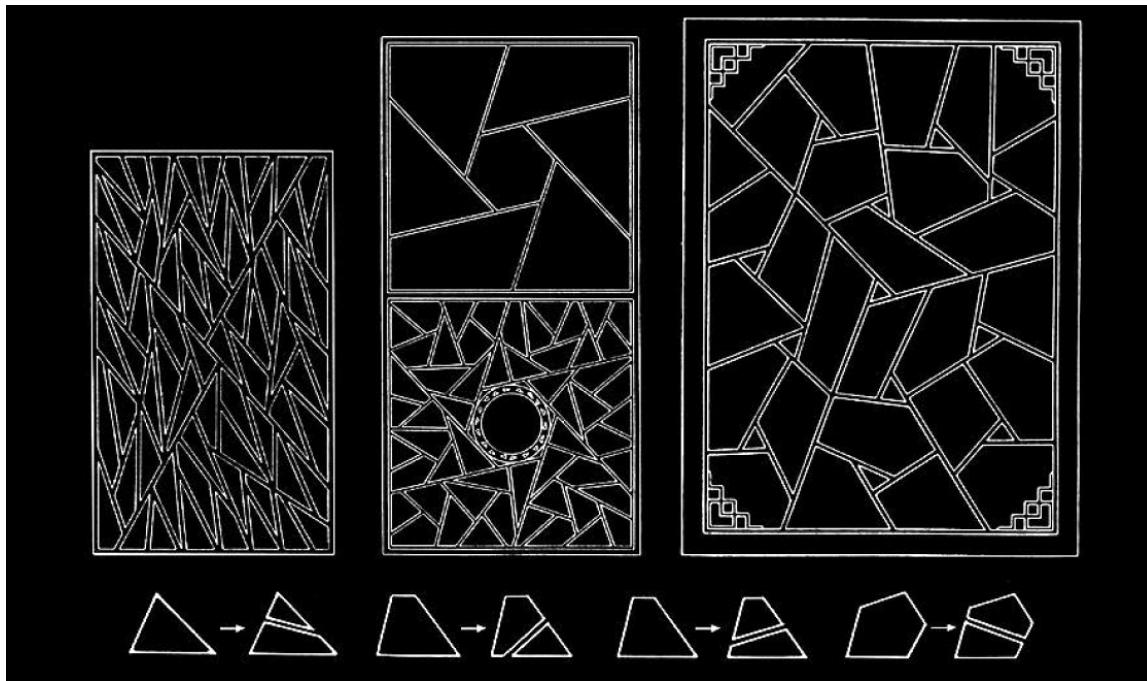


Figure 3.1 Prusinkiewicz' s L-system interpreter concept.

Courtesy of Stelios Manousakis. Used with permission. From "Musical L-Systems." Master's Thesis, Royal Conservatory, The Hague, 2006.

- States determine intervals, not absolute values
- Suggest application to other parameters: tempo, amplitude, and position of sound in space
- Creative applications in the visual arts and architecture

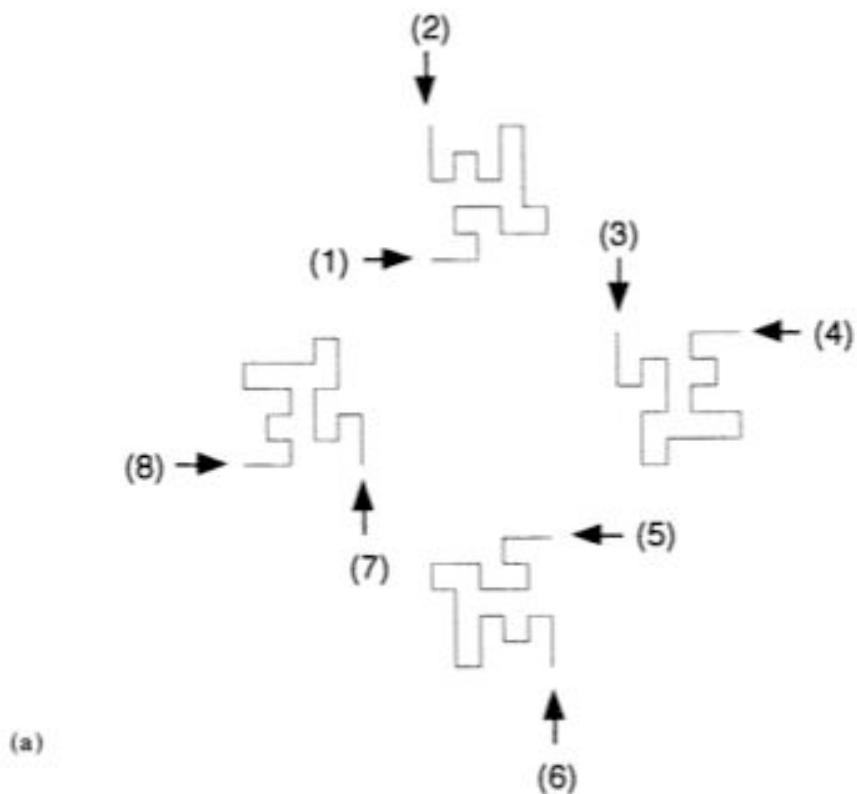
Stiny, G. and J. Gips. 1972. "Shape Grammars and the Generative Specification of Painting and Sculpture." In *Information Processing 71*. C. V. Freiman, ed. Amsterdam: North Holland. 1460-1465. Internet: <http://www.shapegrammar.org/ifip/>.



Courtesy of George Stiny. Used with permission.

### 19.11. Reading: Mason and Saffle

- Mason, S. and M. Saffle. 1994. "L-Systems, Melodies and Musical Structure." *Leonardo Music Journal* 4: 31-38.
- Are deterministic CA always fractal?
- The basic mapping (after Prusinkiewicz)



(1) Original



(2) Original



(3) 90 degree clockwise rotation



(4) 90 degree clockwise rotation



(5) 180 degree rotation



(6) 180 degree rotation



(7) 270 degree clockwise rotation



(8) 270 degree clockwise rotation



(b)

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 Source: Mason, S. and M. Saffle. "L-Systems, Melodies and Musical Structure." *Leonardo Music Journal* 4 (1994): 31-38.

- What are some alternative ways the 2D turtle graphics can be mapped and as musical values?
- Is it significant that “any melody can be modeled with an L-system, including the songs of aboriginal hunters, the plainchants of the medieval christian liturgy, the themes of beethoven’s symphonies and popular song tunes,” as the authors claim?
- What is the implied connection between fractals and beauty. Is this connection sufficiently supported?

## 19.12. A Grammar Specification String and Python Implementation

- Define a grammar in two required parts: alphabet and rules

Both are specified in key{value} pairs

Rules are specified as source{destination} pairs

```
a{3}b{-2} @ a{b} b{a}
```

- Optionally include the axiom (one chosen at random otherwise)

```
a{3}b{-2} @ a{b} b{a} @ baba
```

- Permit one to many rules of any size

```
a{3}b{-2} @ a{ba} b{abb} @ baba
```

- Permit context sensitivity as many to one or many to many rules (not yet implemented)

```
a{3}b{-2} @ aa{ba} bab{abb} @ baba
```

- Match any source as pattern specified with quasi regular expressions (not yet implemented)

```
a{3}b{-2} @ *aa{ba} b*b{abb} bb*{abb} @ baba
```

- Configure non-deterministic destinations as two or more weighted options

Weights can be specified with a floating point or integer value following the destination

```
a{3}b{-2} @ a{ba|ab} b{abb=3|aa=2} @ baba
```

- Can create a grammar instance and view step-wise output

```
>>> from athenaCL.libATH import grammar
>>> g = grammar.Grammar()

>>> g.load('a{3}b{-2} @ a{b} b{a} @ baba')
>>> g.next(); g.getState()
'baba'
>>> g.next(); g.getState()
'abab'
>>> g.next(); g.getState()
'baba'
```

```

>>> g.load('a{3}b{-2} @ a{ba|ab} b{abb=3|aa=2} @ baba')
>>> g.next(); g.getState()
'abbababbba'
>>> g.next(); g.getState()
'baaaabbababbbaabbabbabaaba'
>>> g.next(); g.getState()
'aaabbaabbaabbaabaabaaaaabbbaababbabbabaabbba'

```

- Can translate grammar string back into a list of source values

```

>>> from athenaCL.libATH import grammar
>>> g = grammar.Grammar()
>>> g.load('a{3}b{-2} @ a{ba|ab} b{abb=3|aa=2} @ baba')
>>> g.next(); g.getState()
'abbababbba'
>>> g.getState(values=True)
[3.0, 3.0, 3.0, -2.0, 3.0, -2.0, -2.0, 3.0]

```

### 19.13. Grammar as ParameterObject

- The grammarTerminus ParameterObject

```

:: tpv grammar
Generator ParameterObject
{name,documentation}
grammarTerminus      grammarTerminus, grammarString, stepCount, selectionString
Description: Produces values from a one-dimensional string
rewrite rule, or generative grammar. The terminus, or final
result of the number of generations of values specified by
the stepCount parameter, is used to produce a list of
defined values. Values are chosen from this list using the
selector specified by the selectionString argument.
Arguments: (1) name, (2) grammarString, (3) stepCount, (4)
selectionString {"randomChoice", "randomWalk",
"randomPermutate", "orderedCyclic",
"orderedCyclicRetrograde", "orderedOscillate"}

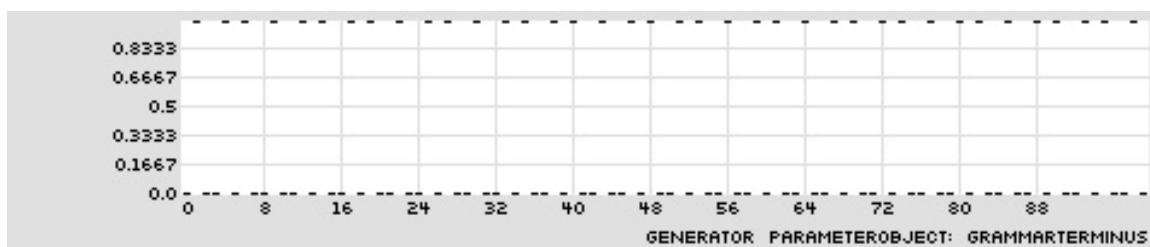
```

- The Lindenmayer alga model after 10 generations

```

:: tpmap 100 gt,a{0}b{1}@a{ab}b{a}@b,10,oc
grammarTerminus, a{0}b{1}@a{ab}b{a}@b, 10, orderedCyclic
TPmap display complete.

```



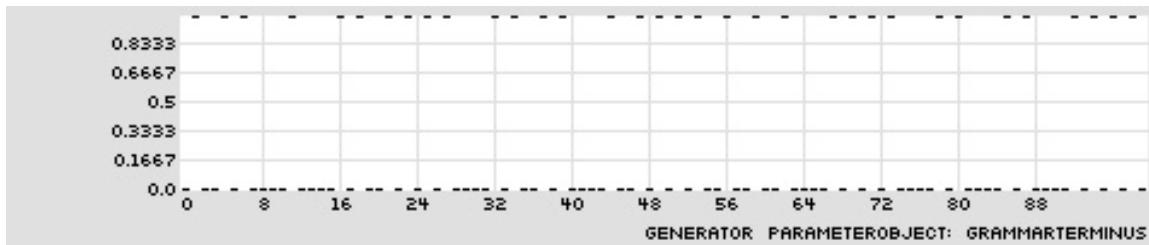
- Modified Lindenmayer alga model after 10 generations with non-deterministic rule variation

```

:: tpmap 100 gt,a{0}b{1}@a{ab}b{a|aaa}@b,10,oc

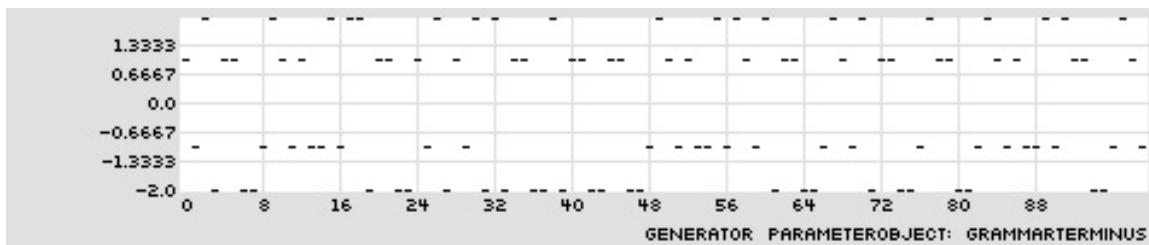
```

```
grammarTerminus, a{0}b{1}@a{ab}b{a=1|aaa=1}@b, 10, orderedCyclic
TPmap display complete.
```



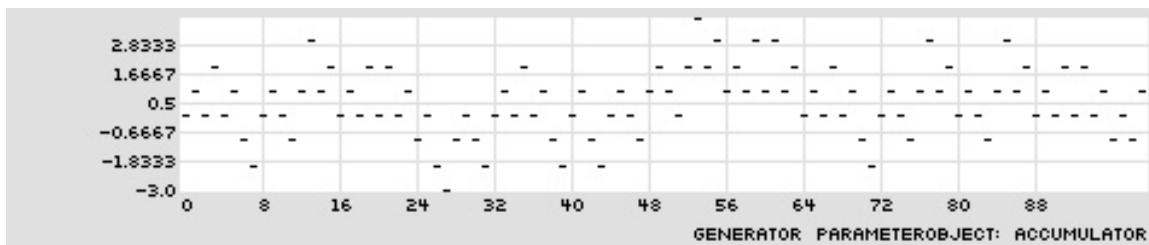
- Four state deterministic grammar

```
:: tpmap 100 gt,a{1}b{-1}c{2}d{-2}@a{ab}b{cd}c{aadd}d{bc}@ac,10,oc
grammarTerminus, a{1}b{-1}c{2}d{-2}@a{ab}c{aadd}b{cd}d{bc}@ac, 10,
orderedCyclic
TPmap display complete.
```



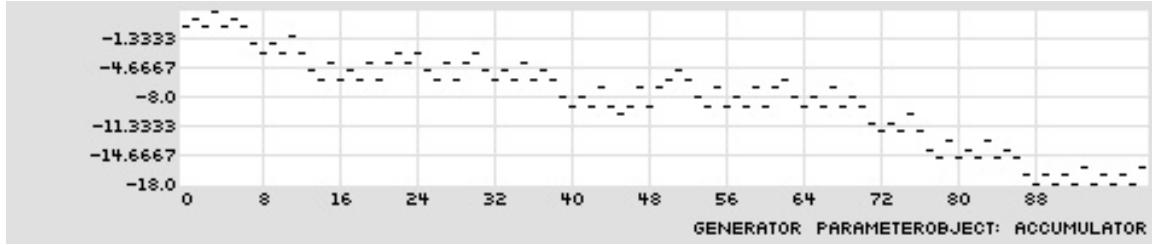
- Four state deterministic grammar placed in an Accumulator PO

```
:: tpmap 100 a,0,(gt,a{1}b{-1}c{2}d{-2}@a{ab}b{cd}c{ad}d{bc}@ac,10,oc)
accumulator, 0, (grammarTerminus, a{1}b{-1}c{2}d{-2}@a{ab}c{ad}b{cd}d{bc}@ac,
10, orderedCyclic)
TPmap display complete.
```



- Four state non-deterministic grammar placed in an Accumulator PO

```
:: tpmap 100 a,0,(gt,a{1}b{-1}c{2}d{-2}@a{ab}b{cd}c{ab|ca}d{ba|db}@ac,10,oc)
accumulator, 0, (grammarTerminus,
a{1}b{-1}c{2}d{-2}@a{ab}c{ab=1|ca=1}b{cd}d{ba=1|db=1}@ac, 10, orderedCyclic)
TPmap display complete.
```



- Alternative approaches to PO interface?
- Mappings and applications in athenaCL?

## 19.14. Grammar States as Accent Patterns

- Can treat the grammar alphabet as parameter values: integers, floating point values
  - Command sequence:
    - emo mp
    - tmo lg
    - tin a 60
    - *non deterministic binary algae generator applied to accent*
- tie r pt,(c,8),(c,1),(gt,a{0}b{1}@a{ab}b{a|aaa}@b,10,oc)
- tie a c,1
  - *four state deterministic applied to pulse multiplier*
- tie r pt,(c,8), (gt,a{1}b{2}c{4}d{8}@a{ab}b{cd}c{aadd}d{bc}@ac,10,oc),(c,1)
- *four state deterministic applied to amplitude with different start string*
- tie a gt,a{.25}b{.5}c{.75}d{1}@a{ab}b{cd}c{aadd}d{bc}@bbc,6,oc
- *four state deterministic applied to transposition with different start string*
- tie f gt,a{0}b{1}c{2}d{3}@a{ab}b{cd}c{aadd}d{bc}@dc,6,oc
- *four state non-deterministic applied to transposition with different start string*
- tie f gt,a{0}b{1}c{2}d{3}@a{ab}b{cd|aa}c{aadd|cb}d{bc|a}@dc,6,oc
- eln; elh

## 19.15. Grammar States as Pitch Values

- Can treat the grammar alphabet as specific pitch values
- Command sequence:
  - `emo m`
  - `tmo lg`
  - `tin a 32`
  - *four state deterministic applied to pulse multiplier*  
`tie r pt,(c,8), (gt,a{1}b{2}c{4}d{8}@a{ab}b{cd}c{aadd}d{bc}@ac,8,oc),(c,1)`
  - `tie o c,-2`
  - *four state deterministic applied to transposition with different start string*  
`tie f gt,a{0}b{7}c{8}d{2}@a{ab}b{cd}c{aadd}d{bc}@ad,6,oc`
  - *four state deterministic applied to amplitude with different start string*  
`tie a gt,a{.25}b{.5}c{.75}d{1}@a{ab}b{cd}c{aadd}d{bc}@bbc,6,oc`
  - `eln; elh`

## 19.16. Grammar States as Pitch Transpositions

- Can treat the grammar alphabet as transpositions iteratively processed through an Accumulator
- Command sequence:
  - `emo m`
  - `tmo lg`
  - `tin a 15`
  - *four state deterministic applied to pulse multiplier*  
`tie r pt,(c,8), (gt,a{1}b{2}c{4}d{8}@a{ab}b{cd}c{aadd}d{bc}@ac,8,oc),(c,1)`
  - *four state deterministic applied to accumulated transposition with different start string*  
`tie f a,0,(gt,a{1}b{-1}c{7}d{-7}@a{ab}b{cd}c{ad}d{bc}@ac,10,oc)`
  - *four state deterministic applied to amplitude with different start string*

tie a gt,a{.25}b{.5}c{.75}d{1}@a{ab}b{cd}c{aadd}d{bc}@bbc,6,oc

- eln; elh

## 19.17. Grammar States as Path Index Values

- Can treat the grammar alphabet as index values from the Path iteratively processed through an Accumulator
- Command sequence:

- emo m

- *create a single, large Multiset using a sieve*

pin a 5@0|7@2,c2,c7

- tmo ha

- tin a 6

- *constant rhythm*

tie r pt,(c,4),(c,1),(c,1)

- *select only Multiset 0*

tie d0 c,0

- *select pitches from Multiset using accumulated deterministic grammar starting at 12*

tie d1 a,12,(gt,a{1}b{-1}c{2}d{-2}@a{ab}b{cd}c{ad}d{bc}@ac,10,oc)

- *create only 1 simultaneity from each multiset; create only 1-element simultaneities*

tie d2 c,1; tie d3 c,1

- *four state deterministic applied to amplitude with different start string*

tie a gt,a{.25}b{.5}c{.75}d{1}@a{ab}b{cd}c{aadd}d{bc}@bbc,6,oc

- eln; elh

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